

Approximating Large Matrix A

1. **Importance and Motivation.** Using SVD we can get the approximation of A in the following way.

Let A be our matrix, defined as a set of column vectors. Establish G as a set of Gaussian random row vectors. Define their product Y .

$$Y = AG$$

$$= \begin{bmatrix} a_1 & \dots & a_n \end{bmatrix} \begin{bmatrix} g_1 \\ \vdots \\ g_n \end{bmatrix}$$

Now derive Q from QR decomposition, $Q = \text{qr}(Y, 0)$. We can define $B = Q^*A$, with decomposition $B = \hat{U}DV^*$, and we claim that

$$QQ^*A \approx A$$

$$QB \approx A$$

$$Q - \text{svd}(B) \approx A$$

This approximation of A can be widely used in a ton of applications, *not just SVD!* Therefore it's important to understand its derivation and usefulness.

2. **Properties and Definition of Q .** We have shown that $QQ^*A \approx A$, and we can examine the properties of QQ^* . We claim that QQ^* is actually a projector onto the range of Q .

To be a projector, the following properties must be satisfied.

1. $\hat{P}^2 = \hat{P}$
2. $\hat{P}\vec{v} = \vec{v}$ for some vector \vec{v} .

These can be trivially shown using the definition of \hat{P} , which implies the following two statements are equivalent.

- $\text{range}(A) \subseteq \text{range}(Q)$
- $A = QQ^*A$.

In a sense, Q captures the range of A .

We can see that these properties hold as long as $\underbrace{QQ^*A}_B \approx A$ and B is smaller than A .

3. **Determining Q .** We've defined Q to be the result of QR decomposition on our Y matrix (our random sampling of A), but we have to be careful how we define that random sampling.

$$Y = \begin{matrix} A & G \\ m \times n & n \times (k+p) \end{matrix}$$

Much of the algorithm is now determined by $k+p$.

- If $k+p$ is small, A isn't sampled enough and our approximation is no longer accurate.
- If $k+p$ is large, B grows large and we've lost the reason why we're trying to approximate A in the first place.

So how do we determine Q ?

4. Adaptive Range Finder. ¹

Given $m \times n$ matrix A , a tolerance ϵ , and an integer r , find Q such that

$$\|(I - QQ^*)A\| \leq \epsilon$$

holds with probability at least $1 - \min\{m, n\} 10^{-r}$.

Data: A, ϵ, r

Result: Q

Draw standard Gaussian vectors $w^{(1)}, \dots, w^{(r)}$ of length n .

for $i = 1, 2, \dots, r$ **do**

$y^{(i)} = Aw^{(i)}$

end

$j = 0$

$Q^{(0)} = []$;

/ the $m \times 0$ empty matrix */*

while $\max\{\|y^{(j+1)}\|, \|y^{(j+2)}\|, \dots, \|y^{(j+r)}\|\} > \epsilon/10\sqrt{2/\pi}$ **do**

$j += 1$

 Overwrite $y^{(j)} = \left(I - Q^{(j-1)}(Q^{(j-1)})^*\right)y^{(j)}$

$q^{(j)} = y^{(j)} / \|y^{(j)}\|$

$Q^{(j)} = [Q^{(j-1)} q^{(j)}]$

 Draw a standard Gaussian Vector $w^{(j+r)}$ of length n .

$y^{(j+r)} = \left(I - Q^{(j-1)}(Q^{(j-1)})^*\right)Aw^{(j+r)}$

for $i = (j + 1) : (j + r - 1)$ **do**

 Overwrite $y^{(i)} = y^{(i)} - q^{(j)}\langle q^{(j)}, y^{(i)} \rangle$

end

end

$Q = Q^{(j)}$

¹Algorithm 4.2 in Halko, Martinsson, Tropp, Page 25, <http://arxiv.org/pdf/0909.4061.pdf>